

Type derivation for extended example

Steffen Zschaler

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This document is an appendix to the paper “Towards Constraint-Based Model Types: A Generalised Formal Foundation for Model Genericity” showing the type derivation of the second example. This type derivation (see Fig. 1) shows that we can then also pass in instances of **Mandatory-Start**. We use the following abbreviations:

$$\begin{aligned}
 Cons_2 &= \neg isAbstract(StateMachine) \wedge \neg isAbstract(State) \\
 Cons'_2 &= Cons_2 \wedge \forall e : (StateMachine, -, -) : lower(e) = 0 \\
 Cons'^*_2 &= Cons_2 \wedge \forall e : (State, -, -) : lower(e) = 0 \\
 Cons''_2 &= Cons'^*_2 \vee \forall e : (-, initialState, State) : lower(e) \leq 1 \\
 T_2 &= (\{StateMachine, State\}, \\
 &\quad \{(StateMachine, initialState, State)\}, \\
 &\quad Cons_2) \\
 T'_2 &= (\{StateMachine, State\}, \\
 &\quad \{(StateMachine, initialState, State)\}, \\
 &\quad Cons'_2) \\
 T'^*_2 &= (\{StateMachine, State\}, \\
 &\quad \{(StateMachine, initialState, State)\}, \\
 &\quad Cons'^*_2) \\
 T''_2 &= (\{StateMachine, State\}, \\
 &\quad \{(StateMachine, initialState, State)\}, \\
 &\quad Cons''_2) \\
 d_2 &= [s \mapsto (\top, T_2)] \\
 d'_2 &= d_2 [sm \mapsto (StateMachine, T'_2)] \\
 d''_2 &= d'_2 [sm \mapsto (StateMachine, T''_2)] \\
 d'''_2 &= d''_2 [s \mapsto (\top, T''_2)] \\
 d''''_2 &= [s \mapsto (\top, T''_2)]
 \end{aligned}$$

$$\begin{array}{c}
\frac{\text{StateMachine} \in \{\text{StateMachine}, \text{State}\} \quad \text{Cons}_2 \Rightarrow \neg \text{isAbstract}(\text{StateMachine})}{d_2 \vdash \text{StateMachine} : \{c : \text{Class} \mid \neg \text{isAbstract}(c)\}} \text{CLASS} \\
\frac{\frac{d_2 \vdash \text{StateMachine} : \{c : \text{Class} \mid \neg \text{isAbstract}(c)\}}{d_2 \vdash \text{new StateMachine}() : (\text{StateMachine}, T_2')} \text{NEW} \quad \text{Var}}{\vdash \text{sm} := \text{new StateMachine}() : d_2 \rightarrow d_2'} \text{VAR} \\
\frac{\frac{\frac{\frac{\frac{\text{StateMachine} \in \{\text{StateMachine}, \text{State}\} \quad \text{Cons}_2 \Rightarrow \neg \text{isAbstract}(\text{StateMachine})}{d_2' \vdash \text{sm} : (\text{StateMachine}, T_2')} \text{CTX1} \quad \text{canNavigate}(\text{StateMachine}, T_2', \text{initialState}, \text{State})}{\vdash \text{sm.initialState} += \text{new State}() : d_2' \rightarrow d_2''} \text{APND2} \quad \text{NEW}}{\vdash \text{sm.initialState} += \text{new State}() ; s += \text{sm} : d_2' \rightarrow d_2'''} \text{SEQ} \quad \text{APND1}}{\vdash \text{sm := new StateMachine}() ; \text{sm.initialState} += \text{new State}() ; s += \text{sm} : d_2 \rightarrow d_2'''} \text{SUB} \\
\vdash \text{sm := new StateMachine}() ; \text{sm.initialState} += \text{new State}() ; s += \text{sm} : d_2 \rightarrow d_2'''} \text{SUB} \\
\vdash \text{sm := new StateMachine}() ; \text{sm.initialState} += \text{new State}() ; s += \text{sm} : d_2'' \rightarrow d_2''' \text{OPDEF} \\
\vdash \text{smo NewSM}(s, T_2') \{ \text{sm} := \text{new StateMachine}() ; \text{sm.initialState} += \text{new State}() ; s += \text{sm} \} : T_2' \text{OPDEF}
\end{array}$$

Figure 1: Type derivation for an extended version of the NewSM program